

csp2hc - Rules

Marcel Oliveira¹ and Jim Woodcock²

¹ Departamento de Informática e Matemática Aplicada, UFRN, Brazil

² Department of Computer Science, University of York, UK

1 Global Environments, Functions, Variables and Constants

– Environments

- Processes information
 - * *ProcArgsEnv*
 - Type: *String* \leftrightarrow seq *Argument*
 - Description: Maps processes names to the names of their arguments
 - * *ProcDefsEnv*
 - Type: *String* \leftrightarrow *ProcBody*
 - Description: Maps processes names to their bodies
 - * *MainProc*
 - Type: *ProcBody*
 - Description: declares the main behaviour of the system
- Types information
 - * *TypeBranchesEnv*
 - Type: *String* \leftrightarrow seq *String*
 - Description: Maps type names to constants that represents the values that variables of each type may assume. For example, if we have the two datatypes in the CSP specification

$$\begin{aligned} \text{datatype } LETTER &= A \mid B \mid C \\ \text{datatype } ID &= \text{Alpha}.LETTER \mid \text{Unknown} \end{aligned}$$

we get the following *TypeBranchesEnv* environment

$$\begin{aligned} LETTER &\mapsto \langle C, B, A \rangle \\ ID &\mapsto \langle \text{Unknown}, \text{Alpha}_C, \text{Alpha}_B, \text{Alpha}_A \rangle \end{aligned}$$

- * *TagsParentEnv*
 - Type: *String* \leftrightarrow *String*
 - Description: Maps constructor (tag) names to the name of the type whose declaration used it. For the example above, we have the following *TagsParentEnv* environment

$$\text{Alpha} \mapsto ID$$

- * *TagTypesEnv*
 - Type: $String \rightarrow \text{seq } String$
 - Description: Maps constructor (tag) names to a sequence of types. For the example above, we have the following *TagsParentEnv* environment

$$Alpha \mapsto \langle LETTER \rangle$$

- * *TagValuesEnv*
 - Type: $String \rightarrow \text{seq } String$
 - Description: Maps constructor (tag) names to a list of list of possible values for each type used in its declaration.

$$Alpha \mapsto \langle \langle C, B, A \rangle \rangle$$

- * *ArgTypeEnv*
 - Type: $String \rightarrow (String \rightarrow String)$
 - Description: For each process name, returns a function that maps the name of the process arguments to its Handel-C type.
- Constants and functions information
 - * *ConstFunArgsEnv*
 - Type: $String \rightarrow \text{seq } String$
 - Description: Maps the function names into the list of its arguments names (constants names are mapped into an empty list)
- Channels information
 - * *CTypeEnv*
 - Type: $String \rightarrow \text{seq } String$
 - Description: Maps channel names into the list of types used in their respective declarations
 - * *CCommEnv*
 - Type: $\text{seq } String$
 - Description: List containing the names of the channels that communicate values using either ! or ?
 - * *CProjEnv*
 - Type: $\text{seq } String$
 - Description: List containing the names of the channels that communicate values using projections $c.*$
 - * *CInOut*
 - Type: $String \rightarrow String \rightarrow \{IN, OUT\}$
 - Description: For each process name, returns a function that yields the usage of each channel name.
 - Example: If a channel c is an input channel within process P , then we have that $\{P \mapsto \{c \mapsto IN\}\}$ is a member of *CInOut*.

– Functions

- *Bits(n)*:
 - * Description: returns the number of bits needed to represent n values.
- *Bitwise(total, on)*

- * Description: returns a string with *total* characters, in which all characters are 0, but the *on*-th character is 1. If *on* is negative, all the characters are 0.

- * Example:

$$\begin{aligned} \text{Bitwise}(4, 3) &= 1000 \\ \text{Bitwise}(4, 0) &= 0001 \\ \text{Bitwise}(4, -1) &= 0000 \end{aligned}$$

- $\Pi(lis)$
 - * Description: distributed cartesian product of a list of lists.
- $\text{OrderedIntRange}(n)$
 - * Description: returns a list of integers that can be represented using *n* bits ordered by their representation as unsigned integers.
 - * Example:

$$\text{OrderedIntRange}(4) = \langle 0, 1, 2, 3, 4, 5, 6, 7, -8, -7, -6, -5, -4, -3, -2, -1 \rangle$$

- $\text{IntName}(n)$
 - * Description: returns the name of the constant that represents an integer *n*.
 - * Example:

$$\begin{aligned} \text{IntName}(0) &= \text{integer_0} \\ \text{IntName}(-1) &= \text{integer_neg_1} \end{aligned}$$

- $\text{Separate}(sep, lst)$
 - * Description: returns the string containing the strings in *lst* separated by *sep*.
 - * Example:

$$\text{Separate}(|, \langle C, B, A \rangle) = C \mid B \mid A$$

- $\text{HCType}(csptype)$
 - * Description: returns the string that declares the Handel-C type that corresponds to the given CSP_M type.
 - * Example:

$$\text{HCType}(T) = \begin{cases} \text{integer} & \text{if } T = \text{Int} \\ \text{boolean} & \text{if } T = \text{Bool} \\ T & \text{otherwise} \end{cases}$$

– Variables

- $\text{IntBits} : \text{int}$: number of bits used to represent the integers
- $\text{MutualRec} : \text{bool}$: indicates if the specification presents mutual recursion or not
- $\text{TailRec} : \mathbb{P} \text{String}$: Set containing the name of the processes that are tail recursive.

– Constants

- ϵ : the empty string

2 Translation Rules

2.1 Programs

$\llbracket Prog \rrbracket^{CSPProg} =$
 `declareClockAndSync()` (Page 4)

$\llbracket datatype\ boolean = true \mid false$
 `datatype integer = (IntBranches(IntBits))`
 $Prog \rrbracket^{Types}$ (Page 5)

$\llbracket Prog \rrbracket^{ConsAndFun}$ (Page 5)

`declareChannels()` (Page 6)

`declareAuxMacros()` (Page 4)

$\llbracket Prog \rrbracket^{Proc}$ (Page 7)

$\llbracket MainProc \rrbracket^{Main}$ (Page 11)

$IntBranches(n) = Separate(1, (IntName) \mapsto (OrderedIntRange(n)))$

2.2 Clock and Sync Values

```
declareClockAndSync() =  
    set clock = external "clock1";  
    typedef unsigned int 1 SYNC;  
    const SYNC syncout = 0;  
    SYNC syncin;
```

2.3 Auxiliary Macros

```
declareAuxMacros() =  
    unsigned 32 random_var =1;  
    macro proc random(INTEGER_I){  
        INTEGER_I =(INTEGER_I<-21)@  
                (INTEGER_I[12]^INTEGER_I[30])@  
                (INTEGER_I[10:1]^INTEGER_I[29:20]);  
    macro expr IS_EMPTY_SET(SET_S) = (SET_S == 0);  
    macro expr SET_UNION(SET_S,SET_T) = (SET_S | SET_T);  
    macro expr SET_DIFF(SET_S,SET_T) = (SET_S & (~SET_T));  
    macro expr SET_INTER(SET_S,SET_T) = (SET_S & SET_T);
```

2.4 Types

$$\begin{aligned} \llbracket \text{datatype } T = T_decl \text{ Prog} \rrbracket^{Types} = & \\ & \text{let } branches = TypeBranchesEnv(T_decl) \\ & \text{in } \left(\begin{array}{l} \text{if } T \neq \text{integer} \\ \text{then } declareType(T, branches) \\ \text{else } \epsilon \end{array} \right) \\ & declareTypeSets(T, branches) \\ & declareLookupTable(T, TagsParentEnvs \triangleright T) \\ \llbracket Prog \rrbracket^{Types} & \\ \llbracket P \text{ Prog} \rrbracket^{Types} = \llbracket Prog \rrbracket^{Types} & \\ \\ declareType(T, \langle b_0, \dots, b_n \rangle) = & \\ & \#define T unsigned int Bits(n) \\ & \#define b_n 0 \\ & \dots \\ & \#define b_0 n \\ & \#define T_card (n + 1) \\ \\ declareTypeSets(T, \langle b_0, \dots, b_n \rangle) = & \\ & \#define T_set unsigned int T_card \\ & \#define b_n_set 0bBitwise(n + 1, n) \\ & \dots \\ & \#define b_0_set 0bBitwise(n + 1, 0) \\ & \#define T_set_nil 0bBitwise(n + 1, -1) \\ & static T_set T_set_LUT[T_card] = {b_n_set, ..., b_0_set}; \\ \\ declareLookupTables(T, \langle \rangle) = \epsilon & \\ declareLookupTables(T, tag_0 : tags) = & \\ & \text{let } \langle T_0, \dots, T_m \rangle = TagTypesEnv(tag_0), \\ & \quad Values = \langle \langle T_0_v_0, \dots \rangle, \dots, \langle T_m_v_0, \dots \rangle \rangle = TagValuesEnv(tag_0) \\ & \text{let } \langle \langle T_0_v_0, \dots, T_m_v_0 \rangle, \dots \rangle = \Pi Values \\ & \text{in static T T_tag_0_LUT}[T_0_card] \dots [T_m_card] = \{ \\ & \quad \{ \dots \{ T_0_v_0, \dots, T_m_v_0 \}, \dots \} \\ & \quad \} \\ & declareLookupTables(T, tags) \end{aligned}$$

2.5 Constants and Functions

$$\begin{aligned} \llbracket Prog \rrbracket^{ConsAndFun} = & \\ & declareCFPrototypes \\ & \llbracket Prog \rrbracket^{CFTrans} \end{aligned}$$

```

declareCFPrototypes() =
  let ConstFunArgsEnv = {c1 ↦ args1, ..., cn ↦ argsn}
  in macro expr c1(Separate(, , args1));
  ...
  macro expr cn(Separate(, , argsn));

[[C = body Prog]]CFTrans =
  if C ∈ dom ConstFunArgsEnv
  then let args = ConstFunArgsEnv(C)
        in macro expr name(Separate(, , args)) = ( [[body]]Exp );
  else ε
  [[Prog]]ConsAndFun

[[F(args) = body Prog]]CFTrans =
  let args = ConstFunArgsEnv(C)
  in macro expr name(Separate(, , args)) = ( [[body]]Exp ); (Page 11)
  [[Prog]]ConsAndFun

[[P Prog]]ConsAndFun = [[Prog]]ConsAndFun

```

2.6 Channels

```

[[Prog]]Channels =
  chan SYNC INEXISTENT_CHANNEL; declareChannels()
  declareChannelsCons()

declareChannels() =
  let {c1 ↦ types1, ..., cn ↦ typesn} = CTypeEnv
  in declareChannel(c1)
  ...
  declareChannel(cn)

declareChannel(c) = chan (declareChannelType(c)) c (declareChannelArray(c));

declareChannelType(c) =
  let types = CTypeEnv(c)
      arrayDim = GetArrayDim(c, #types)
  in if (arrayDim == #types)
      then SYNC
      else HCTYPE(last(types))

getArrayDim(c) =
  let size = #CTypeEnv(c)
  in if c ∈ CCommEnv
      then size - 1
      else size

```

```

declareChannelArray(c) =
  if c ∈ CProjEnv
  then let ⟨T0, ..., Tn⟩ = CTypeEnv(c)
        arrayDim = GetArrayDim(c, #types)
        in if (arrayDim == n + 1)
            then [HCTYPE(T0)_card] ... [HCTYPE(Tn-1)_card] [HCTYPE(Tn)_card]
            else [HCTYPE(T0)_card] ... [HCTYPE(Tn-1)_card]
  else ϵ

```

```

declareChannelCons() =
  let {c1, ..., cn} = dom CTypeEnv
  in [[datatype CHAN = chan_c_1 | ... | chan_c_n]]Types

```

2.7 Processes

```

[[Prog]]Proc =
  if mutualRec
  then declareMutualRecCons()
        [[Prog]]ProcDecl
  else declareProcPrototypes()
        [[Prog]]ProcDeclMR

```

```

declareMutualRecCons() = let {P0, ..., Pn} = dom ProcDefsEnv
                        in #define P_n n
                        ...
                        #define P_0 0

```

```

declareProcPrototypes() = let {P0, ..., Pn} = dom ProcDefsEnv
                        in inline void P_0(declareFormalArgs(P0));
                        ...
                        inline void P_n(declareFormalArgs(Pn));

```

```

declareFormalArgs(P) =
  let ⟨arg0, ..., argn⟩ = ProcArgsEnv(P)
  in ( (ArgTypeEnv P arg0) arg_0, ... , (ArgTypeEnv P argn) arg_n )

```

```

 $\llbracket P(\text{args}) = \text{body } Prog \rrbracket^{ProcDecl} =$ 
  if  $P \in \text{dom } ProcDefsEnv$ 
  then inline void  $P(\text{declareFormalArgs}(P))\{$ 
    if  $P \in TailRec$ 
    then boolean KEEP_LOOPING;
      KEEP_LOOPING = true;
      while(KEEP_LOOPING){
        KEEP_LOOPING = false;
         $\llbracket PDef \rrbracket^{PBody} P$ 
      }
    else  $\llbracket PDef \rrbracket^{PBody} P$ 
  }
   $\llbracket Prog \rrbracket^{ProcDecl}$ 
else  $\llbracket Prog \rrbracket^{ProcDecl}$ 

```

```

 $\llbracket Prog \rrbracket^{ProcDeclMR} =$ 
  let  $nproc = \#ProcDefsEnv$ 
     $\{P_0, \dots, P_n\} = \text{dom } ProcDefsEnv$ 
  in declareFormalArgsCopies( $P_0$ )
  ...
  declareFormalArgsCopies( $P_n$ )

  inline void MUTUAL_REC(int  $Bits(nproc)$  PROGRAM_COUNTER){
    boolean KEEP_LOOPING;
    KEEP_LOOPING = true;
    while(KEEP_LOOPING){
      KEEP_LOOPING = false;
      switch(PROGRAM_COUNTER){
        let  $\langle P_0, \dots, P_n \rangle = \text{dom } ProcDefsEnv$ 
        in case  $P\_0$  : {  $\llbracket ProcDefsEnv(P_0) \rrbracket^{PBody} P_0$  }
        ...
        case  $P\_n$  : {  $\llbracket ProcDefsEnv(P_n) \rrbracket^{PBody} P_n$  }
        default: KEEP_LOOPING = false;
      }
    }
  }

```

```

declareFormalArgsCopies( $P$ ) =
  let  $\langle arg_0, \dots, arg_n \rangle = ProcArgsEnv(P)$ 
  in ( $ArgTypeEnv P arg_0$ )  $P\_local\_arg\_0$ ;
  ...
  ( $ArgTypeEnv P arg_n$ )  $P\_local\_arg\_n$ ;

```


Process Body

$$\llbracket SKIP \rrbracket^{PBody} N = \epsilon$$

$$\llbracket STOP \rrbracket^{PBody} N = \text{INEXISTENT_CHANNEL?syncin};$$

$$\llbracket P1; P2 \rrbracket^{PBody} N = (\llbracket P1 \rrbracket^{PBody} N); (\llbracket P2 \rrbracket^{PBody} N)$$

$$\begin{aligned} \llbracket \text{if } C \text{ then } P1 \text{ else } P2 \rrbracket^{PBody} N = \\ \text{if } (\llbracket C \rrbracket^{Exp}) \{ \llbracket P1 \rrbracket^{PBody} N \} \text{ else } \{ \llbracket P2 \rrbracket^{PBody} N \} \end{aligned}$$

$$\begin{aligned} \llbracket P1 \mid \sim \mid P2 \rrbracket^{PBody} N = \\ \text{random}(\text{random_var}); \\ \text{if } ((\text{random_var} \% 2) == 0) \{ \\ \llbracket P1 \rrbracket^{PBody} N \} \\ \text{else} \{ \\ \llbracket P2 \rrbracket^{PBody} N \} \end{aligned}$$

$$\begin{aligned} \llbracket C \&P \rrbracket^{PBody} N = \\ \text{if } (\llbracket C \rrbracket^{Exp}) \{ \llbracket P \rrbracket^{PBody} N \} \text{ else } \{ \llbracket STOP \rrbracket^{PBody} N \} \end{aligned}$$

Prefix

$$\begin{aligned} \llbracket comm \rightarrow P \rrbracket^{PBody} N = \\ \text{seq} \{ \llbracket comm \rrbracket^{DeclInVar} \llbracket comm \rrbracket^{Comm}; \llbracket P \rrbracket^{PBody} N \} \end{aligned}$$

$$\begin{aligned} \llbracket c \text{ proj } ?x \rrbracket^{DeclInVar} &= H\text{Ctype}(\text{last}(C\text{TypeEnv}(c))) \text{ x}; \\ \llbracket comm \rrbracket^{DeclInVar} &= \epsilon \end{aligned}$$

$$\begin{aligned} \llbracket c \text{ proj} \rrbracket^{Comm} N &= \text{if } N \in \text{dom } C\text{InOut} \wedge c \in \text{dom } C\text{InOut}(N) \\ &\quad \text{then if } C\text{InOut } N \text{ } c = IN \\ &\quad \quad \text{then } c(\llbracket proj \rrbracket^{Proj} (C\text{TypeEnv } c)) \text{?syncin} \\ &\quad \quad \text{else } c(\llbracket proj \rrbracket^{Proj} (C\text{TypeEnv } c)) \text{!syncout} \end{aligned}$$

$$\begin{aligned} \llbracket c \text{ proj } ?x \rrbracket^{Comm} &= c(\llbracket proj \rrbracket^{Proj} (C\text{TypeEnv } c)) \text{?x} \\ \llbracket c \text{ proj } !e \rrbracket^{Comm} &= c(\llbracket proj \rrbracket^{Proj} (C\text{TypeEnv } c)) \text{!} \llbracket e \rrbracket^{Exp} \end{aligned}$$

$$\begin{aligned} \llbracket .e_0 \dots .e_n \rrbracket^{Proj} \langle T_0, \dots, T_n \rangle = \\ \llbracket (\text{if } T_0 = \text{Int then (unsigned) else } \epsilon) \llbracket e_0 \rrbracket^{Exp} \rrbracket \\ \dots \\ \llbracket (\text{if } T_n = \text{Int then (unsigned) else } \epsilon) \llbracket e_n \rrbracket^{Exp} \rrbracket \end{aligned}$$

External Choice

$$\begin{aligned}
& \llbracket comm_0 \rightarrow P_0 \square \dots \square comm_n \rightarrow P_n \rrbracket^{PBody} N = \\
& \quad \llbracket comm_0 \rrbracket^{DeclInVar} \\
& \quad \dots \\
& \quad \llbracket comm_1 \rrbracket^{DeclInVar} \\
& \quad \text{case } \llbracket comm_0 \rrbracket^{Comm} : \{ \llbracket P_0 \rrbracket^{PBody} N \}; \text{ break;} \\
& \quad \dots \\
& \quad \text{case } \llbracket comm_n \rrbracket^{Comm} : \{ \llbracket P_0 \rrbracket^{PBody} N \}; \text{ break;}
\end{aligned}$$

Parallel Composition

$$\begin{aligned}
& \llbracket P_1 [\mid cs \mid] P_2 \rrbracket^{PBody} N = \\
& \quad TranslateConc(N, P_1, P_2, ((\alpha(P_1) \cap \alpha(P_2)) \setminus cs)) \\
& \llbracket P_1 [cs_1 \mid \mid cs_2] P_2 \rrbracket^{PBody} N = \\
& \quad TranslateConc(N, P_1, P_2, ((\alpha(P_1) \cap \alpha(P_2)) \setminus (cs_1 \cap cs_2))) \\
& \llbracket P_1 \mid \mid \mid P_2 \rrbracket^{PBody} N = \\
& \quad TranslateConc(N, P_1, P_2, (\alpha(P_1) \cap \alpha(P_2)))
\end{aligned}$$

$$\begin{aligned}
& TranslateConc(N, P_1, P_2, \emptyset) = \\
& \quad \text{par } \{ \{ \llbracket P_1 \rrbracket^{PBody} N \} ; \{ \llbracket P_1 \rrbracket^{PBody} N \} \}
\end{aligned}$$

Process/Function Call

$$\begin{aligned}
& \llbracket P(e_0, \dots, e_n) \rrbracket^{PBody} N = \\
& \quad \text{if } P \in \text{dom } ProcArgsEnv \\
& \quad \text{then let } \langle arg_0, \dots, arg_n \rangle = ProcArgsEnv(P) \\
& \quad \text{in if } (MutualRec \vee P = N) \\
& \quad \quad \text{then } P_local_arg_0 = \llbracket e_0 \rrbracket^{Exp} ; \\
& \quad \quad \dots \\
& \quad \quad P_local_arg_n = \llbracket e_n \rrbracket^{Exp} ; \\
& \quad \quad \left(\begin{array}{l} \text{if } (MutualRec) \\ \text{then if } N = \text{mainp} \\ \quad \text{then } MUTUAL_REC(P); \\ \quad \text{else } PROGRAM_COUNTER = P; \\ \text{else } \epsilon \end{array} \right) \\
& \quad \quad \text{KEEP_LOOPING} = \text{true}; \\
& \quad \quad \text{else } P(\llbracket e_0 \rrbracket^{Exp}, \dots, \llbracket e_n \rrbracket^{Exp}); \\
& \quad \text{else } P(\llbracket e_0 \rrbracket^{Exp}, \dots, \llbracket e_n \rrbracket^{Exp});
\end{aligned}$$

2.8 Main Process

$$\llbracket main \rrbracket^{Main} =$$

```
void main(){
     $\llbracket main \rrbracket^{PBody}$  mainp
}
```

2.9 Expressions

The function $\llbracket exp \rrbracket^{Exp}$ returns the corresponding Handel-C code for a given CSP_M expression.

3 Extensions

3.1 Allowing interleaved events

- Declaration of the branches
- Changes to translation (`this.useSemaphors`)